Hallucination-Resilient LLM-Driven Sound and Tunable Static Analysis

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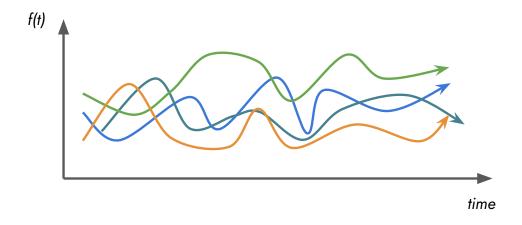




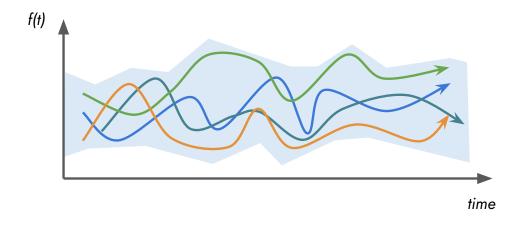




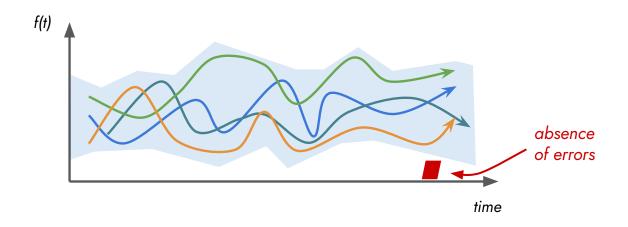




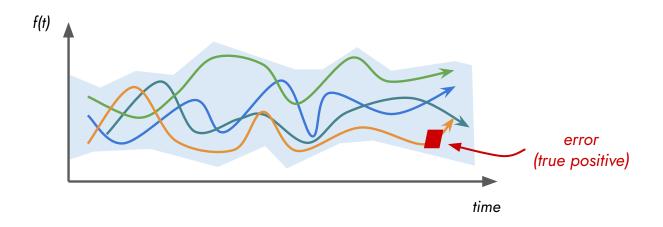
concrete traces



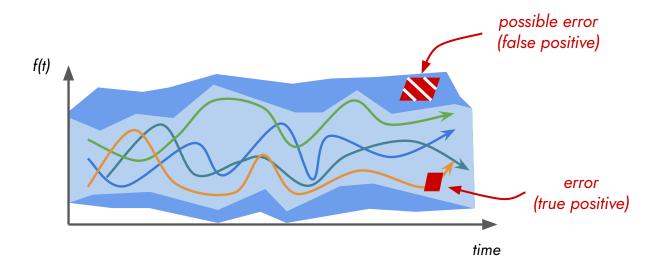
soundness: over-approximations



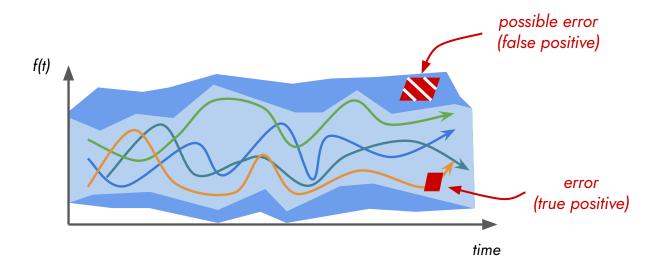
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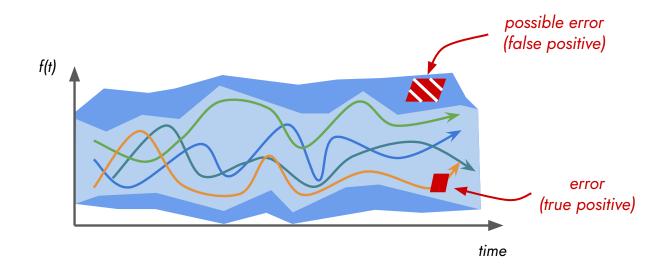
soundness: over-approximations



coarser-grained over-approximations



Goals
Soundness: detect all true errors
Precision: reduce false errors



So, you want to use LLMs for static analysis, but how?

Prompting LLMs

```
J demo.java 1 X
Users > xiangqian > Downloads > J demo.java > ...
       import java.util.ArrayList;
      public class Demo {
          public static int foo(int a, int b) {
               if (Math.abs(b) > 1)
                  System.out.println(a / b);
               return b;
          Run | Debug
          public static void main(String[] args) {
               int x = 1;
               int y = 10 / x;
               ArrayList<Integer> list = new ArrayList<>();
               list.add(x);
               list.add(y);
               list.add(e:0);
              y = x * 1;
               int z = x / y;
              z = list.get(index:2);
              System.out.println(y / z);
               x = foo(y, z);
               System.out.println(x / x);
```

Claude-3.5-Sonnet

Prompting LLMs

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```

```
CHAT

Please mimic the execution of the program and detect divide-by-zero bugs.

J demo.java:1-26
```

prompt

Prompting LLMs

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```
CHAT
                                                                                              prompt
Please mimic the execution of the program and detect divide-by-zero bugs.
   J demo.java:1-26
 8 GitHub Copilot
   Used 1 reference
I'll analyze the code execution step by step and identify potential divide-by-zero bugs.
Analysis
  1. In (m) main():
      o (e) x = 1
      \circ \bowtie y = 10 / x \rightarrow \bowtie y = 10 (safe, x = 1)

    Create ArrayList and add: [1, 10, 0]

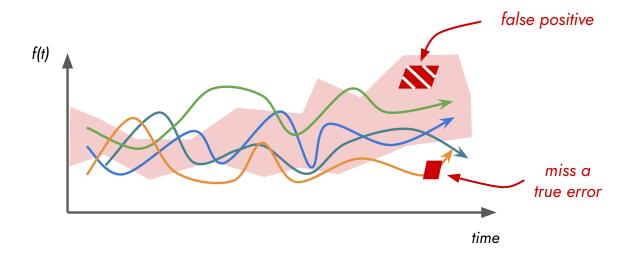
      reasoning &
      \circ \bowtie z = x/y <math>\rightarrow \bowtie z = 1 (safe, y = 1)
                                                                                        analysis result
      \circ \bowtie z = list.get(2) \rightarrow \bowtie z = 0
      • System.out.println(y/z) \rightarrow DIVIDE BY ZERO ERROR! (y = 1, z = 0)
      o Program would crash here, but if it continued:
      \circ \bowtie x = foo(y, z) with \bowtie y = 1, z = 0

    Inside (a) foo(): doesn't print since (a) Math.abs(0) ≤ 1

      o Returns o z which is 0
      ○ System.out.println(x / x) \rightarrow DIVIDE BY ZERO ERROR! (x = 0)
Divide-by-Zero Bugs Found:
  1. Line 21: System.out.println(y / z) - when so z = 0
  2. Line 24: System.out.println(x / x) - when x = 0
```

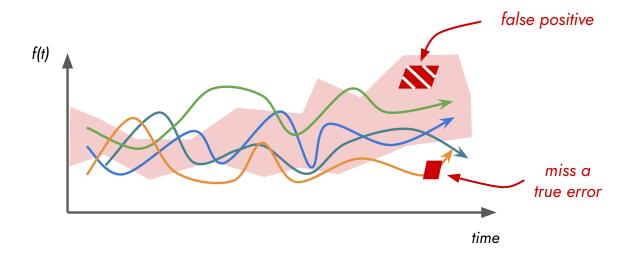
Claude-3.5-Sonnet

Can we trust LLM-reported analysis results?



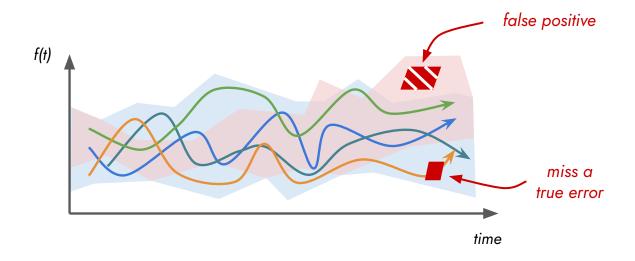
LLMs are untrustworthy because of hallucination.

Can we trust LLM-reported analysis results?



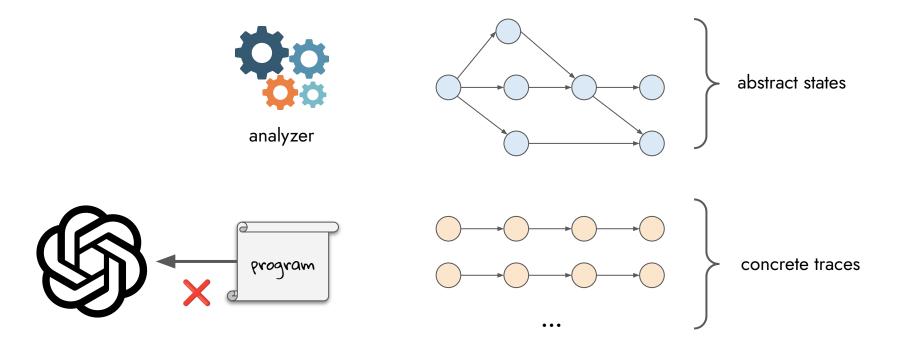
Can we validate the analysis results from LLMs?

Can we trust LLM-reported analysis results?

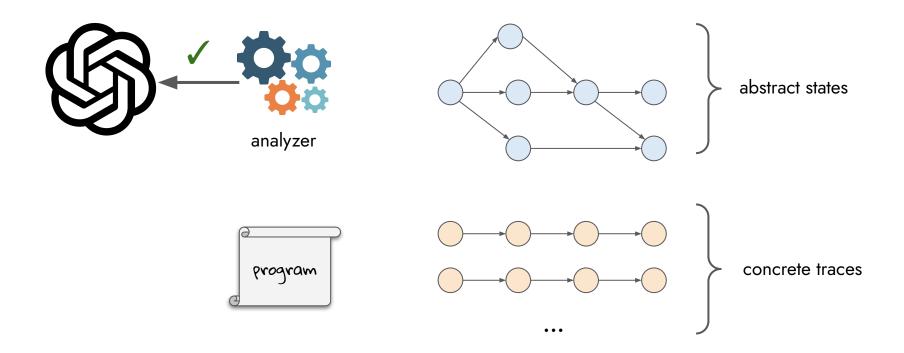


Can we validate the analysis results from LLMs? Yes, but this becomes a static analysis problem...

Instead of analyzing target programs, use LLMs to "analyze" the analysis!



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Recipe:

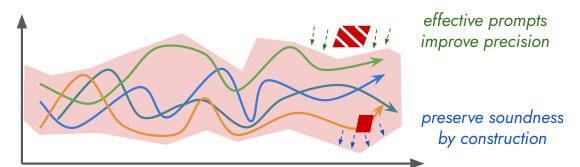
- Identify precision-critical decisions in the static analysis
 - Any decision does not hinder soundness,
 - but only impacts precision
- During the analysis, use LLMs as the decision maker
 - Prompt describes the process/implementation of static analysis
 - Prompt contains current abstract state and other (meta-) information
 - Prompt instructs to decide for better precision

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- Control flow analysis:
 - o determines the function to be called at each call-site
 - undecidable problem once data- and control-flow are entangled (e.g. when with first-class functions in Scheme, dynamic dispatch in Java, indirect call with pointers in C, etc.)
- Abstracting abstract machine (AAM) framework [ICFP '10]
 - o analysis is defined as an abstraction of a concrete abstract machine (e.g. CEK machine)

$$\Sigma \in \widehat{\text{State}} \triangleq (\text{Expr} + \widehat{\text{Value}}) \times \widehat{\text{Env}} \times \widehat{\text{BStore}} \times \widehat{\text{KStore}} \times \widehat{\text{Kont}} \times \widehat{\text{Time}}$$

o result is an over-approximated transition graph of abstract states

• [Recipe step 1] Identify precision-critical decisions in the static analysis

• [Recipe step 1] Identify precision-critical decisions in the static analysis

```
\begin{array}{l} \textit{tick} &: \widehat{\mathsf{State}} \to \widehat{\mathsf{Time}} \\ \\ \textit{alloc}_{\nu} : \widehat{\mathsf{State}} \times \mathsf{Var} \times \widehat{\mathsf{Time}} \to \widehat{\mathsf{BAddr}} \\ \\ \textit{alloc}_{\kappa} : \widehat{\mathsf{State}} \times \widehat{\mathsf{Expr}} \times \widehat{\mathsf{Env}} \times \widehat{\mathsf{BStore}} \times \widehat{\mathsf{Time}} \to \widehat{\mathsf{KAddr}} \end{array}
```

LLM as abstract address allocator

• [Recipe step 1] Identify precision-critical decisions in the static analysis

 $tick : \widehat{State} \to \widehat{Time}$

 $alloc_{v}: \widehat{State} \times Var \times \widehat{Time} \rightarrow \widehat{BAddr}$

 $alloc_{\kappa}: \widehat{\mathsf{State}} \times \mathsf{Expr} \times \widehat{\mathsf{Env}} \times \widehat{\mathsf{BStore}} \times \widehat{\mathsf{Time}} \to \widehat{\mathsf{KAddr}}$

LLM as abstract address allocator

Soundness-safe:

Any allocation strategy from a finite set is sound (a Posteriori Soundness Theorem)

Precision-critical:

Addresses decide the degree of approximation (e.g. context-sensitivity, ...)

• [Recipe step 2] During the analysis, use LLMs as the decision maker.

```
Your task is to analyze the process of a small-step abstract interpreter that transforms an abstract machine state into possible successor states.
  I will provide you the data structures used in the abstract interpretation first
 It will provide you the data strangures used in the addrest interpretation rirst,

Ithen give your the current analysis state. Ask you to analyze it, and ask you to decide the context-sensitivity or
 |
|The "State" field is an abstract state, which can be see of the following:
 |enum State:
| case Estate(e: Expr. p: Env. e-| Estare, p.| Estore, k: Eant, t: Time:
  The "Expr" of "State" field is the current program/expression being evaluated, which is defined by the following algebraic data type:
 | case Lit(n: Int)
  case transplay String, angl Expr)
case Mindplep: String, That Expr, that Expr)
case Var(x: String)
case Lam(x: String, body: Expr)
   case App(f) Expr. argi Expr)
 | case Letrec(s: String, rhs: Eapr, body: Eapr)
| case Degin(eaprs: List(Eapr))
|The input program is already renamed so that variable names are unique.
 |
|The "Value" of "State" field is the value being evaluated, which is defined by the following algebraic data type:
 The "Env" field is the environment mapping variable names to their abstract addresses BAddr.
[type Env = Map[String, BAddr]
[case class BAddrix: String, instrumentation: List[Any])
 The "KStore" field is the store mapping abstract binding addresses to sets of continuations, each map is represented with ->-
 |
|The "Mont"/"continuation" field is the current continuation, which is defined by the following algebraic data type:
  case KillaryEp(op) String, or thy, ky KAGEr)
   Case KArgie: Exp., p: Env., k: Akder)
Case KYun(lam: Expr.Lam, p: Env., k: KAder)
Case KLetix: String, p: Env., body: Expr., k: KAder)
Case KLetrec(x: String, xa: Bader, p: Env., body: Expr., k: KAder)
   case KRegin(exprs: List(Expr), X: KAddr)
 |The "Time" field is a list of program call sites expression that represent the history of
  calls leading to the current state (context sensitivity).
|Type Time = List(Expr)
 Your input is a JSGN object with the following schema:
                                       // for BASST, WASST, and "TLCK" (ALMKES HAVE THES FIELD)
   "time": Time.
   "time": Time, // for RAddr, RAdd

"Source-mapression": Expr., // for RAddr celly

"target-expression": Expr., // for RAddr celly

"target-expression": Expr., // for RAddr celly

"target-incise-store": Store, // for RAddr celly
 .

You should analyze the current "state" and return an abstract address for better analysis $(goal).
  since if the binding address already exists, reusing it decreases the precision of the analysis.
   "reason: strang "geory-type" immediate or "TLCK", "yarlable": String, // For Baddr only "time" Book, // For both Baddr and Kaddr "source-expression": Book, // For Kaddr only
   "target-expression": Bool, // for Käddr only
"target-environment": Bool, // for Käddr only
"target-sinding-store": Bool, // for Käddr only
   "k": Int represented as String, // for fick only
 Pield "reason" should explain why you chose this address for better s(goal) by analyzing the current state.
 The reases you give should be specific to the current state and should not be generic.
[Field "query-type" should be either "BAdde", "KAdde", or "Tick" according to the "query-type" field in the imput query.
IYOU may set time to false if you think it is not needed for the analysis.
 [Field "variable" should the same variable name as in the input query for binding addresses.
(DIT) of Temporyper in Table?

Final Tempor-depositor Thouse or treat of the continuation powers (should be instrumented with the Secret expression, elemented false, Final Tempor-deposition) though the treat of the continuation powers (should be interrupted expression) who do be true of the continuation powers and the continuation powers (should be interrupted expression) and the continuation powers (should be interrupted expression).
|Field "target-binding-store" should be true if the continuation address should be instrumented with the target binding store, otherwise false.
(Final "" Books in the name (represented as a quoted strong) of content-neutrity beets to use in the analysis. 
Glowes the returned value of "", the majoris will the scholl-content-neutrity.

[I.L. in code the current call sepression is presence to the history, and then the that first K call sites of the whole history), reposably, from "is a sealth water, thus of or. I the "returned to the history".
```

 [Recipe step 2] During the analysis, use LLMs as the decision maker.

```
The second section of the second section of the second second second second section se
```

```
|The "State" field is an abstract state, which can be one of the following: |enum State: | case EState(e: Expr, \rho: Env, \sigma_v: BStore, \sigma_k: KStore, k: Kont, t: Time) | case VState(v: Value, \rho: Env, \sigma_v: BStore, \sigma_k: KStore, k: Kont, t: Time) | case ErrState()
```

Describe data structures used in the analyzer

```
Interpripage "Swellow of "Nation" for "Natio
```

• [Recipe step 2] During the analysis, use LLMs as the decision maker.

```
|Your input is a JSON object with the following schema:
1{
  "state": State,
  "query-type": "BAddr" or "KAddr" or "Tick",
  "variable": String.
                                       // for BAddr only
  "time": Time,
                                       // for BAddr, KAddr, and "Tick"
  "source-expression": Expr,
                                       // for KAddr only
  "target-expression": Expr,
                                       // for KAddr only
  "target-environment": Env,
                                       // for KAddr only
  "target-binding-store": BStore,
                                       // for KAddr only
```

Describe input JSON format for each query.

```
Your task is to analyze the process of a small-step abstract interpreter that transforms an abstract machine state into possible successor state
   case estate(e) bugs, o) boy, e-| Botore, e,| Motore, N| Mont, t| Time
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The input program is already renamed so that variable names are unique.
 The "Value" of "State" field is the value being evaluated, which is defined by the following algebraic data type:
|type Env = Map(String, BAddr)
|case class BAddrix: String, instrumentation: List[Any])
   case Kinaryto(so: String, o: Env. h: KAGOri
   case KESTORK(sp: String, rhs: Expr, p: Env, k: KAder)
case KESTOR((sp: String, ths: Value, k: KAder)
case KERTOR((sp: Expr, p: Env, k: KAder)
   case KFue(Last Expr.Las, p: Env, k: KAddr)
case KLetix: String, p: Env, body: Expr. k: KAddr)
case KLetrec(x: String, xa: Bäddr, p: Env, body: Expr. k: KAddr)
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        time": Time.
                                                                                            // for Baddr, Waddr, and "Tick" (ALMRYS HAVE THES FIELD
 You should analyze the current "state" and return an abstract address for better analysis $(goal)
     "target-expression": Bool, // for Kiddr only
"target-environment": Bool, // for Kiddr only
"target-sinding-store": Bool, // for Kiddr only
      "k": Int represented as String, // for Tick only
 The reason you give should be specific to the current state and should not be generic.
Field "query-type" should be either "Badde", "MAdde", or "Tick" according to the "query-type" field in the input query.
You may set time to false if you think it is not needed for the analysis.
 Field "variable" should the same variable name as in the input query for binding addresses.
 Only if "super-type" is "ABDET;

The Contraction Super-type is "ABDET;

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The Contraction Super-type is super-type in the contraction Super-type is super-type in the Super-type in the Contraction Super-type is super-type in the Contraction Super-type in the Contraction Super-type in the Contraction Super-type in the Contraction Super-type in Super-type Supe
|Field "target-binding-store" should be true if the continuation address should be instrumented with the target binding store, otherwise false.
Transally, field "A" is a small number. The B or L. If you think X by 2 is needed, please justify in the "reason" field
```

• [Recipe step 2] During the analysis, use LLMs as the decision maker.

```
Your output should be only a JSON object with the following schema:
  "reason": string
  "query-type": "BAddr" or "KAddr" or "Tick",,
  "variable": String.
                               // for BAddr only
  "time": Bool,
                                 // for both BAddr and KAddr
  "source-expression": Bool,
                            // for KAddr only
  "target-expression": Bool, // for KAddr only
  "target-environment": Bool, // for KAddr only
  "target-binding-store": Bool, // for KAddr only
  "k": Int represented as String, // for Tick only
```

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case Bindp(op: String, Tha: Expr, tha: Expr)
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     Case Margie: Expr. p: Env. k: Moder)
Case Minister Expr.Lam. p: Env. k: Moder)
Case Minister String, p: Env. body: Expr. k: Moder)
Case Mintrack: String, was Baddr, p: Env. body: Expr. k: Moder)
       case KBeginleagrs: List(Expr), KI KASSY)
     calls leading to the current state (context sensitivity).
 Your input is a 350% object with the following schema
"SIDE" SIDE,

"GROUP SPECT TAKEN" OF "MAKEN" OF "TICK",

"WATERSON'S TOTAGE,

"For Maken, Maken, and "TICK" (ACRES HOR THIS FILE)

"WATERSON'S TIME.

""" STATE SAMEN, MAKEN, AND TICK" (ACRES HOR THIS FILE)
     "Source-expression": Expr. // for Käddr celly
"target-expression": Expr. // for Käddr celly
"target-environment": Env. // for Käddr celly
"target-environment": Env. // for Käddr celly
"target-environment": Expr. // for Käddr celly
                                                                                           // for KAddr only
// for KAddr only
 You may set time to false if you think it is not needed for the analysis.
   Field "variable" should the same variable name as in the input query for binding addresses.
 [DIS] If "Assurtings" is "Made";
If "High "Searc-Seption" hould be form if the continuation powers should be instrumented with the searce expression, intervise Tailor.
If "High "Searc-Seption" hould be true if the continuation approximation and the true of the target expression, therefore Tailor.
If "High "Search-Search Tailor is the continuation appears should be interneted with the target expression, therefore the Tailor "Search-Search Tailor" and Tailor "Search Tailor" and Tailor "Search-Search Tailor" and Tailor "Search Tailor" and Tailor "Search Tailor" and Tailor "Search-Search Tailor" and Tailor "Search Tailor" a
 |Field "target-binding-store" should be true if the continuation address should be instrumented with the target binding store, otherwise false.
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```

• [Recipe step 2] During the analysis, use LLMs as the decision maker.

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```

|You should analyze the current "state" and return an abstract address for better analysis \${goal}.

|You should look at the entries of the existing binding store and continuation store,

|since if the binding address already exists, reusing it decreases the precision of the analysis.

Explain the logic behind allocation and ask for reasoning with the abstract states.

```
The compact is a 20th object with the following schools:

"Suppregnent "Suppress" "Waller or "Title".

"Suppregnent "Suppress" "Waller or "Title".

"Suppregnent "Suppress" "Waller or "Title".

"The Title of Suppress of Sup
```

 A prototype "LLMAAM" implemented in Scala https://github.com/Kraks/llmaam

Preliminary evaluation of micro-benchmarks with GPT-4o-mini

	idid		kcfa2		mj09		loop2	
	#N	#E	#N	#E	#N	#E	#N	#E
0CFA+P4F	42	41	307	306	303	302	110	108
0CFA+AAC	33	32	203	202	179	175	110	108
0CFA+SRC	42	41	246	245	647	644	149	146
1CFA+P4F	33	32	TO	TO	1437	1407	108	106
1CFA+AAC	33	32	336	324	331	323	108	106
1CFA+SRC	42	41	12189	12117	6849	6834	108	106
1000-000	33	32	133	131	199	202	108	106
llmaam	33	32	153	150	268	259	108	106
	33	32	272	265	281	273	108	106

#edges and #states as approximation of precision - smaller is better

Use LLMs to "Analyze" the Analysis!

Recipe for sound, precise, tunable static analysis driven by LLMs:

- Identify precision-critical decisions in static analysis
- During the analysis, use LLMs as meta-analysis and the decision maker

Future opportunities - Analyses:

- Control-flow sensitivity
- Numeric analysis
- Termination analysis
- ...

Future opportunities - Optimizations:

- Prompt optimizations
- Fine-tuning
- Reinforcement learning
- ...